

# Computação Visual e Multimédia

10504: Mestrado em Engenharia Informática

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Chap. 9 — Object Data Structures

## Object Data Structures



## Outline

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- Motivation.
- Geometric structures versus topological structures.
- Topological data structures: introduction.
- Incidence and adjacency relationships.
- Spaghetti data structure.
- DCEL data structure.
- Symmetric data structure.
- Topological inference and reasoning on incidence and adjacency.
- Euler operators (still incomplete!)



## Geometric object data structures

### **Purpose:**

- data structures for representing and manipulating geometric objects in space.

### **Requirement:**

- the stored information must allow for an unambiguous representation of the subdivision.

### **Evaluation:**

- space complexity: amount of space (memory) needed for storing all information (entities and relations) that is explicitly represented
- time complexity of the algorithms for calculating relations that are not explicitly represented



## Topological Data Structures



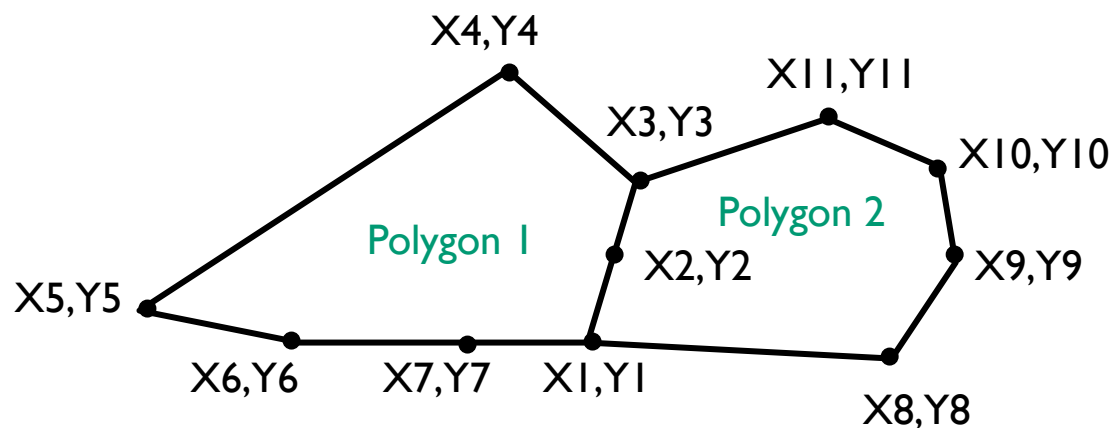
## Topology / connectivity

- Generic sets of entities: vertices, edges and faces
- Overlaid sets of entities: only *meet* and *disjoin*
- **Meet**: topological relation that defines connectivity between entities. Entities of different dimension are “connected” in different ways: relations (vertex-, edge-, face-based)
- **Disjoin**: topological relation that defines the entities of lower dimension are in the boundary of of higher dimension entities.

## Spaghetti data structure

- Spaghetti data structure: represents sets of points, lines and polygons
- Can be used for both generic sets of entities and overlaid sets (plane subdivisions)
- The geometry of any spatial entity is described independently of other entities
- No topology/connectivity information is recorded

- Points, lines and polygons are stored separately.
- For each polygon, we store a (ordered) list of coordinates of points on its boundary.



Polygon 1	Polygon 2
X1,Y1	X8,Y8
X2,Y2	X9,Y9
X3,Y3	X10,Y10
X4,Y4	X11,Y11
X5,Y5	X3,Y3
X6,Y6	X2,Y2
X7,Y7	X1,Y1
X1,Y1	X8,Y8



## Spaghetti data structure: pros & cons

### Advantages:

- simplicity
- easy insertions of new entities (all entities are independent)

### Disadvantages:

- inefficient for topological queries

No easy way of solving queries such as: “do Polygon 1 and 2 share a common bounding line?”  
*Need to analyse all coordinates of points of Polygon 1 and compare with those of Polygon 2 and see if two consecutive pairs are the same: inefficient!!*

- redundancies (and consequently, possible inconsistencies!)

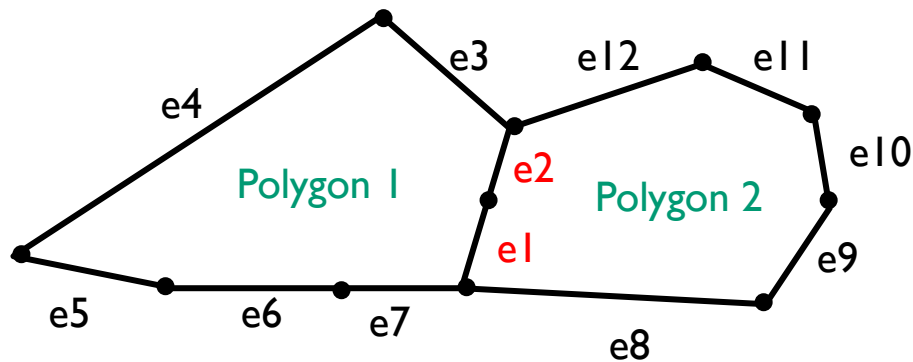
Coordinates of points along common boundary are recorded twice!  
*Redundancy: if we update coordinates of a point, we need to update them everywhere!*

## Topological data structures: motivation

- Storing connectivity information explicitly allows for more efficient spatial *queries*.
- Topology/connectivity: important criterion to establish the *correctness* (integrity, consistency) of geometric objects, with applications in CAD, geographical databases, etc.

### Example:

If we store relation FE explicitly (i.e., for each polygon we store a list of IDs of edges bounding it), the query: “do Polygon 1 and 2 share a common bounding line?” only requires checking whether the two lists contain any common IDs



Polygon 1	Polygon 2
e1	e8
e2	e9
e3	e10
e4	e11
e5	e12
e6	e2
e7	e1

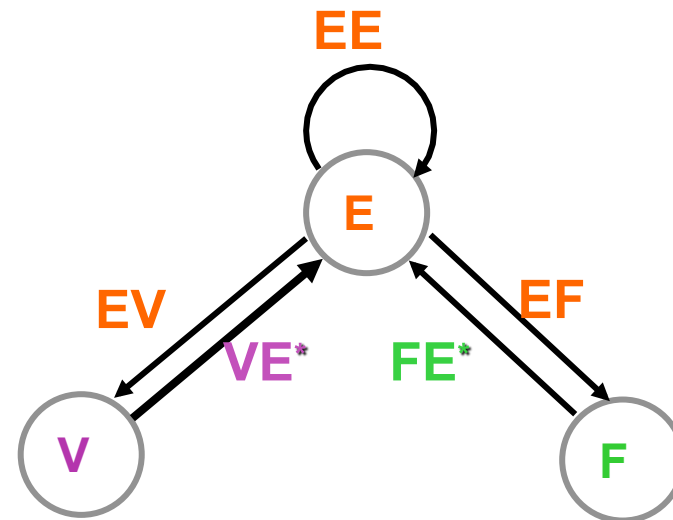


## Doubly-connected edge list (DCEL)

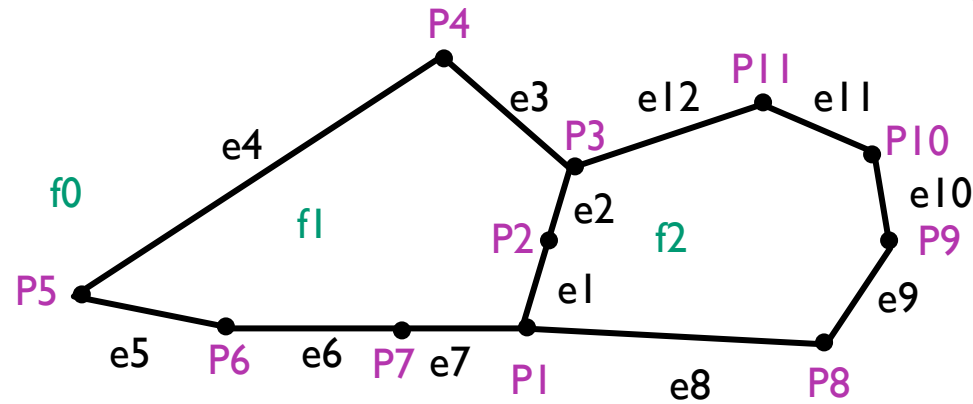
– Preparata and Shamos (1985)

**DCEL** structure stores:

- three sets of entities  $V$ ,  $E$ ,  $F$
- three edge-based relations  $EV$ ,  $EE$ ,  $EF$
- two partial relations:  $FE^*$  and  $VE^*$ 
  - $FE^*$ : associates a face  $f$  with one of the edges bounding  $f$
  - $VE^*$ : associates a vertex  $v$  with one of the edges incident in  $v$



# Example



Entities	
V	P1,P2,..., P11
E	e1, e2,..., e12
F	f0,f1,f2

Edge - based relations			
	EV	EF	EE
e1	P1,P2	f1,f2	e7,e2
e2	P2,P3	f1,f2	e1,e12
e3	P3,P4	f1,f0	e2,e4
e4	P4,P5	f1,f0	e3,e5
e5	P5,P6	f1,f0	e4,e6
e6	P6,P7	f1,f0	e5,e7
e7	P7,P1	f1,f0	e6,e8
e8	P1,P8	f2,f0	e1,e9
e9	P8,P9	f2,f0	e8,e10
e10	P9,P10	f2,f0	e9,e11
e11	P10,P11	f2,f0	e10,e12
e12	P11,P3	f2,f0	e11,e3

Partial relations			
	VE*		FE*
P1	e1	f0	e3
P2	e2	f1	e3
P3	e3	f2	e1
P4	e4		
P5	e5		
P6	e6		
P7	e7		
P8	e8		
P9	e9		
P10	e10		
P11	e11		



## DCEL: space complexity

For every **edge**:

- 3 constant relations are stored (involving 2 entities):  $6e$

For every **face**:

- 1 relation involving one entity:  $f$

For every **vertex**:

- 1 relation involving one entity:  $n$

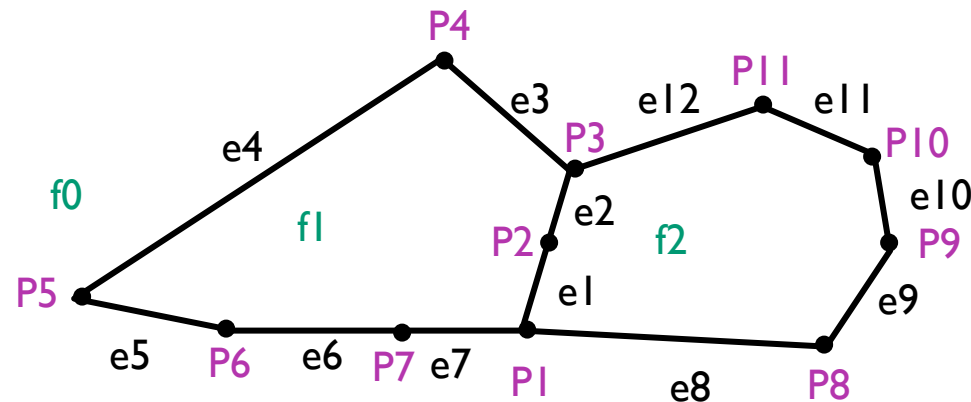
TOTAL space required to represent relations:  $6e + f + n$

For each vertex we also store the two geometric coordinates:  $2n$

## DCEL: time complexity for FE

### Calculating complete relation FE: Obtained by combining FE\* and EE

- For example, given a face  $f_1$ , we find the first bounding edge  $e_3$  using  $FE^*$ . Then using EE we find the successor  $e_4$  of  $e_3$  (in counter-clockwise order) on the boundary of  $f_1$ : if  $e_3$  is oriented in such a way that  $f_1$  is on its left hand side, then  $e_4$  is the second of the two edges associated with  $e_3$  through EE
- We apply the same method (2<sup>nd</sup> element of EE) to obtain all other edges on the boundary of  $f_1$ , until we reach  $e_3$  again.





## DCEL: FV and FF

FV relation:

- FV can be obtained by combining FE and EV: for each bounding edge  $e$  of a given face  $f$  (obtained with FE), we consider its endpoints using EV

FF relation:

- FF can be obtained by combining FE and EF: for each bounding edge  $e$  of a given face  $f$  (obtained with FE), we consider the other face  $f'$  obtained by using EF

VE relation:

- Homework...

VV relation:

- Homework...

VF relation:

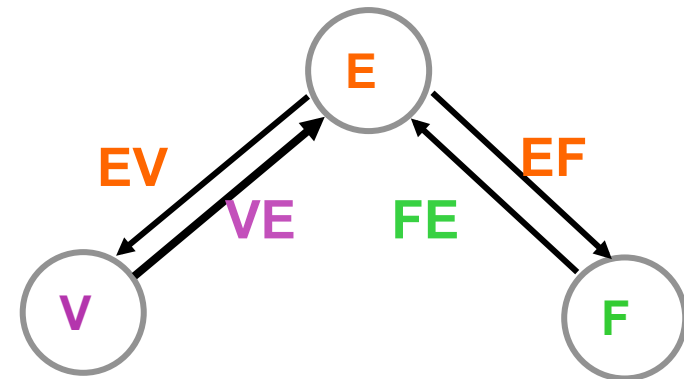
- Homework...

## Symmetric data structure

- Woo (1985)

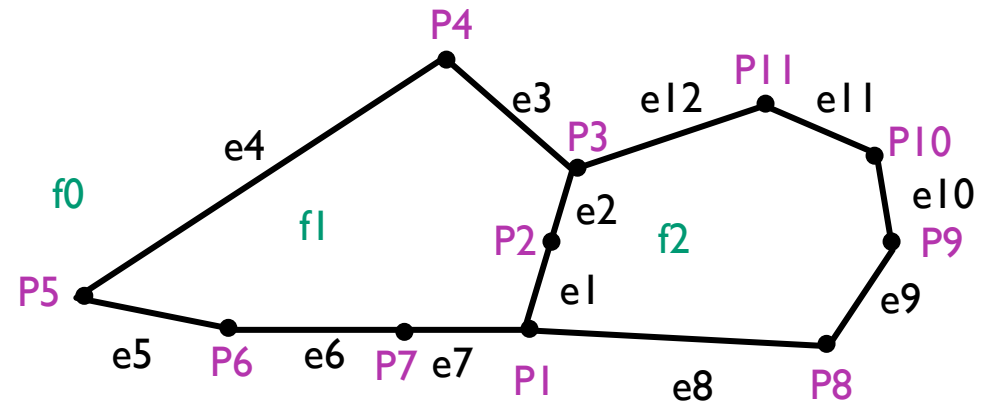
Symmetric structure stores:

- three sets of entities: V, E, F
- relation EV and its inverse VE
- relation FE and its inverse EF



# Example

	<b>EV</b>	<b>EF</b>		<b>VE</b>		<b>FE</b>
e1	P1,P2	f1,f2	P1	e1,e7,e8	f0	e3, e4,e5,e6,e7,e8,e9,e10,e11,e12
e2	P2,P3	f1,f2	P2	e2,e1	f1	e3,e4,e5,e6,e7,e1,e2
e3	P3,P4	f1,f0	P3	e3,e2,e12	f2	e1,e8,e9,e10,e11,e12,e2
e4	P4,P5	f1,f0	P4	e4,e3		
e5	P5,P6	f1,f0	P5	e5,e4		
e6	P6,P7	f1,f0	P6	e6,e5		
e7	P7,P1	f1,f0	P7	e7,e6		
e8	P1,P8	f2,f0	P8	e8,e9		
e9	P8,P9	f2,f0	P9	e9,e10		
e10	P9,P10	f2,f0	P10	e10,e11		
e11	P10,P11	f2,f0	P11	e11,e12		
e12	P11,P3	f2,f0				





## Symmetric structure: space complexity

For every **edge**:

- 2 constant relations are stored (involving 2 entities):  $4e$

For every **face**:

- 1 variable relation (FE). Every edge is common to two faces, so each edge is stored twice:  $2e$

For every **vertex**:

- 1 variable relation (VE). Every edge has two endpoints, so each edge is stored twice:  $2e$

TOTAL space required to represent relations:  $8e$

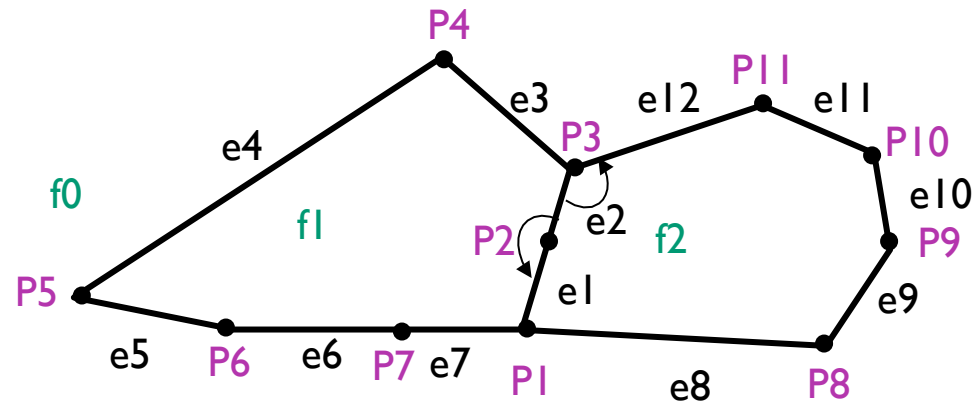
For each vertex we also store the two geometric coordinates:  $2n$



## Symmetric structure: EE

### Calculating relation EE: Obtained by combining EV and VE (or EF and FE)

- For example, if we want to calculate  $EE(e2)=(e1,e12)$ , we retrieve the endpoints  $P2$  and  $P3$  of  $e2$  using EV. To retrieve  $e1$  we consider the successor of  $e2$  in the list associated with  $P2$  through VE (for  $e12$  the successor of  $e2$  in the list associated with  $P3$ ). To do this in constant time, for each edge we need to store the position of the edge in the lists associated to its endpoints through VE.



## Symmetric structure: FF , FV, VV, VF

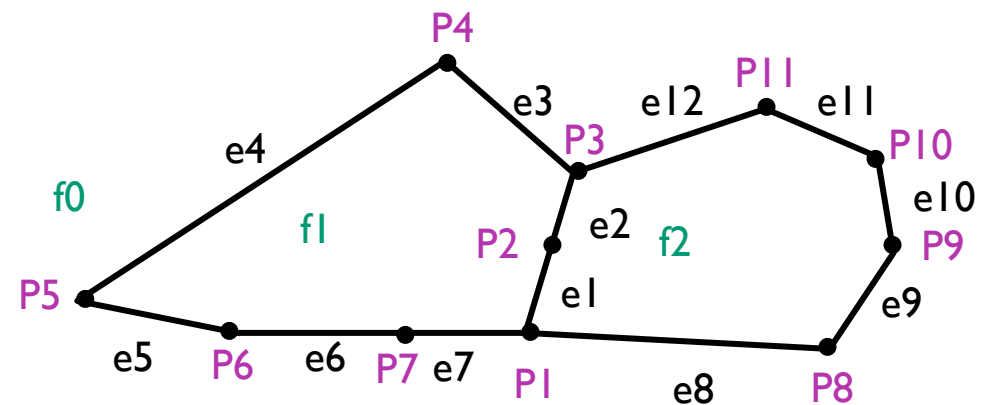
As in DCEL:

- FF: FE+EF
- FV: FE+EV
- VV: VE+EV
- VF: VE+EF

### Example: FF

FF( $f_1$ )=( $f_0, f_2$ ) obtained combining:

- FE( $f_1$ )=( $e_3, e_4, e_5, e_6, e_7, e_1, e_2$ )
- EF( $e_3$ )=( $f_1, f_0$ )
- EF( $e_4$ )=( $f_1, f_0$ )
- EF( $e_5$ )=( $f_1, f_0$ )
- EF( $e_6$ )=( $f_1, f_0$ )
- EF( $e_7$ )=( $f_1, f_0$ )
- EF( $e_1$ )=( $f_1, f_2$ )
- EF( $e_2$ )=( $f_1, f_2$ )





## Euler operators

*Motivation for studying Euler operators:*

- Allow the incremental construction of complex objects from basic building blocks such as vertices, edges and faces.
- Applications: geometric CAD kernels, computational animation systems, etc.



**To be continued....**



## Summary:

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