

Lógica Computacional

LEI, 2014/2015

DI-UBI

Aulas Práticas 24

Dedução Natural em Lógica de Primeira Ordem

Provam-se as seguintes afirmações.

1. $\{\forall x P(x) \vee \forall x Q(x)\} \vdash \forall x (P(x) \vee Q(x))$
2. $\vdash (\forall x P(x) \wedge \forall x Q(x)) \leftrightarrow \forall x (P(x) \wedge Q(x))$
3. $\{\exists x (P(x) \wedge Q(x))\} \vdash \exists x P(x) \wedge \exists x Q(x)$
4. $\vdash (\exists x P(x) \vee \exists x Q(x)) \leftrightarrow \exists x (P(x) \vee Q(x))$
5. $\{\forall x (P(x) \rightarrow Q(x))\} \vdash \forall x P(x) \rightarrow \forall x Q(x)$
6. $\{\exists y \forall x \varphi\} \vdash \forall x \exists y \varphi$
7. $\vdash \exists x \neg P(x) \leftrightarrow \neg \forall x P(x)$
8. $\vdash \forall x \neg P(x) \leftrightarrow \neg \exists x P(x)$
9. $\vdash \exists x \varphi \leftrightarrow \neg \forall x \neg \varphi$
10. $\vdash \forall x \varphi \leftrightarrow \neg \exists x \neg \varphi$
11. $\vdash \forall x \varphi \wedge \psi \leftrightarrow \forall x (\varphi \wedge \psi)$, se $x \notin \text{VL}(\psi)$
12. $\vdash \forall x \varphi \vee \psi \leftrightarrow \forall x (\varphi \vee \psi)$, se $x \notin \text{VL}(\psi)$
13. $\vdash \exists x \varphi \wedge \psi \leftrightarrow \exists x (\varphi \wedge \psi)$, se $x \notin \text{VL}(\psi)$
14. $\vdash \exists x \varphi \vee \psi \leftrightarrow \exists x (\varphi \vee \psi)$, se $x \notin \text{VL}(\psi)$
15. $\vdash \forall x (\psi \rightarrow \varphi) \leftrightarrow \psi \rightarrow \forall x \varphi$, se $x \notin \text{VL}(\psi)$
16. $\vdash \exists x (\psi \rightarrow \varphi) \leftrightarrow \psi \rightarrow \exists x \varphi$, se $x \notin \text{VL}(\psi)$
17. $\vdash \forall x (\varphi \rightarrow \psi) \leftrightarrow \exists x \varphi \rightarrow \psi$, se $x \notin \text{VL}(\psi)$
18. $\vdash \exists x (\varphi \rightarrow \psi) \leftrightarrow \forall x \varphi \rightarrow \psi$, se $x \notin \text{VL}(\psi)$